## Solution of the Exam in Micro A 02/2006

## Question 1

(a) We show that the following preference relation induces the behavior specified in the question:

$$u(x) = \begin{cases} 1 + x_{K+1} & \text{if } x > D\\ \min\left(\frac{x_1}{d_1}, \frac{x_2}{d_2}, \dots, \frac{x_K}{d_K}\right) & \text{otherwise} \end{cases}$$

where  $D = (d_1, d_2..., d_K)$ .

Since  $\min\left(\frac{x_1}{d_1}, \frac{x_2}{d_2}, ..., \frac{x_K}{d_K}\right) \leq 1$  for all bundles in the relevant domain, whenever the consumer can afford more than D he will choose to purchase the bundle D and spend the rest of his income on  $x_{K+1}$ . If he cannot afford more than D then the consumer will behaves according the min function and consume the bundle tD ( $t \leq 1$ ) (where t is a function of his budget).

(b) We will show that behavior does not induce a continuous preference relation (and in particular does not induce a continuous monotonic and convex preference relation).

According to the consumer's behavior  $(d_1, d_2, ..., d_K, \varepsilon) \succ (d_1, d_2, ..., d_K, 0)$  however for every  $\delta > 0$  $(d_1 - \delta, d_2, ..., d_K, \varepsilon) \prec (d_1, d_2, ..., d_K, 0)$  contradicting continuity.

## Question 2

Let  $x \sim_i y$  denote that individual *i* views *x* as "the same as" *y* and  $x \nsim_i y$  denote that individual *i* views *x* as "not the same as" *y*,  $x \sim y$  and  $x \nsim y$  with no index is society's opinion. Let  $(\sim_i)_{i=1,...,N}$  be a profile of equivalence relations then an aggregation method is a function  $F: (\sim_i)_{i=1,...,N} \to E$ .

(a)

- P: For all  $x, y \in X$  and every profile  $(\sim_i)_{i=1,\dots,N}$ , if  $x \sim_i y \ \forall i$  then  $x \sim y$  and if  $x \nsim_i y \ \forall i$  then  $x \nsim y$ .
- I\*: For every  $a, b, c, d \in X$  and any two profiles  $(\sim_i)_{i=1,...,N}$  and  $(\sim'_i)_{i=1,...,N}$  if for all  $i, a \sim_i b$  iff  $c \sim'_i d$  then  $a \sim b$  iff  $c \sim' d$ .

(b)

• Satisfies P but not I\*:  $F((\sim_i)_{i=1,...,N})$  is the most common equivalence relation among  $(\sim_i)_{i=1,...,N}$  (with some tie breaker).

P is satisfied since if all individuals view x as "the same as" y then in particular the most common equivalence relations view x as "the same as" y, thus society views x as "the same as" y (the opposite is true if x is "not the same as" y for all i).

I\* is not satisfied: look at the following example:  $X = \{x, y, z\}$  Let  $(\sim_i)_{i=1,...,5}$  be  $x \sim_i y \sim_i z$  for i = 1, 2,  $x \approx_3 y$  and  $x \sim_3 z$ ,  $x \approx_4 y$  and  $y \sim_4 z$  and  $x \ll_5 y \ll_5 z \ll_5 x$ . The the most common equivalence relation is that of i = 1, 2 so  $x \sim_i y$ . However for  $(\sim'_i)_{i=1,...,5}$  where  $\sim'_i = \sim_i$  for i = 1, 2 and  $\sim'_i = \sim_5$  for i = 3, 4, 5 we have  $x \approx y$ , even though  $x \sim_i y$  iff  $x \sim'_i y$  for all i, contradicting I\*.

• Satisfies I\* but not P:  $F((\sim_i)_{i=1,...,N})$  is  $x \sim y \sim z$  for every profile  $(\sim_i)_{i=1,...,N}$ .

I\* is satisfied since  $F((\sim_i)_{i=1,\ldots,N})$  is constant for every profile.

P is not satisfied since even if  $x \nsim_i y \forall i, x \sim y$  contradicting P.

• Satisfies I\* and P:  $\forall x, y \in X \ x \sim y \text{ iff } x \sim_i y \ \forall i \in G \subset N.$ 

P is satisfied since if all individuals view x as "the same as" y then in particular  $\forall i \in G$  view x as "the same as" y and thus society views x as "the same as" y. If all individuals view x as "not the same as" y then  $\exists i \in G$  that views x as "not the same as" y thus society views x as "not the same as" y.

I\* is satisfied since in any two profiles  $(\sim_i)_{i=1,...,N}$  and  $(\sim'_i)_{i=1,...,N}$  if for all  $i, a \sim_i b$  iff  $c \sim'_i d$  we have  $a \sim_i b$  iff  $c \sim'_i d \forall i \in G$  thus  $a \sim b$  iff  $c \sim' d$ .

(c) We will show that if X includes at least three elements, then the only aggregation method which satisfies P and I<sup>\*</sup> is the aggregation method which determines a subset  $G^* \subseteq N$  such that  $x \sim y$  iff  $x \sim_i y \ \forall i \in D$ .

Define  $\Gamma = \{G \subseteq N \mid \text{for all } x, y \in X, \text{ if for all } i \in G \ x \sim_i y \text{ and for all } j \notin G \ x \approx_j y \text{ then } x \sim y \}$ . Note that if  $G \in \Gamma$  then  $G \neq \emptyset$  since  $G = \emptyset$  would imply that  $x \sim y$  for the profile  $(\sim_i)_{i=1,\ldots,N}$  in which  $x \approx_i y \ \forall i \in N$  contradicting P. Furthermore,  $\Gamma$  is not empty since by P  $N \in \Gamma$ .

If  $G_1, G_2 \in \Gamma$  then  $G_1 \cap G_2 \in \Gamma$ .

We have to show that for any x, y and for any profile  $(\sim'_i)_{i=1,...,N}$  for which  $x \sim'_i y$  for all  $i \in G_1 \cap G_2$ , and  $x \approx'_i y$  for all  $i \in G_1 \cap G_2$  the equivalence relation  $F((\sim'_i)_{i=1,...,N})$  determines that  $x \sim' y$ . By  $I^*$  it is sufficient to show that for one pair a and b, and for one profile  $(\sim_i)_{i=1,...,N}$  that agrees with the profile  $(\sim'_i)_{i=1,...,N}$  on the pair  $\{a, b\}$ , the equivalence relation  $F((\sim_i)_{i=1,...,N})$  determines that  $a \sim b$ . Let  $c \neq a, b$ . Let  $(\sim_i)_{i=1,...,N}$  be a profile satisfying for all  $i \in G_1 \cap G_2$ ,  $a \sim_i b \sim_i c$ , for all  $i \in G_1 \setminus G_2 a \approx_i b$  and  $a \sim_i c$ , for all  $i \in G_2 \setminus G_1$   $a \approx_i b$  and  $b \sim_i c$  and for all  $i \in N \setminus (G_1 \cup G_2)$   $a \approx_i b$  and  $b \approx_i c$ . Since  $G_1 \in \Gamma$   $a \sim c$ . Since  $G_2 \in \Gamma$   $b \sim c$ . By transitivity  $a \sim b$ .

There exists a unique minimal (with respect to inclusion) non empty subset  $G^* \in \Gamma$ .

Assume there does not exist a unique minimal subset  $G^* \in \Gamma$ . Then there exist two subsets  $G_1, G_2 \in \Gamma$  such that  $G_1 \neq G_2$  and  $\nexists G \in \Gamma$  such that  $G \subset G_1$  and  $\nexists G \in \Gamma$  such that  $G \subset G_2$ . By the previous claim  $G_1 \cap G_2 \in \Gamma$  and  $G_1 \cap G_2 \subset G_1$  and  $G_1 \cap G_2 \subset G_1$ . Furthermore,  $G^* \neq \emptyset$  since  $N \in \Gamma$  and  $\emptyset \notin \Gamma$ .

If  $G \in \Gamma$  then for all  $G' \supseteq G$   $G' \in \Gamma$ . Take any  $x, y \in X$ . Let  $(\sim_i)_{i=1,...,N}$  be a profile satisfying for all  $i \in G$   $x \sim_i y \sim_i z$ , for all  $i \in G' \setminus G$   $x \sim_i y$  and  $x \nsim_i z$  and for all  $i \in N \setminus G'$   $x \nsim_i y, y \nsim_i z$ . Since  $G \in \Gamma$  $x \sim z$  and  $y \sim z$ . By transitivity a  $x \sim y$  thus  $G' \in \Gamma$ . This implies that for any profile  $(\sim'_i)_{i=1,...,N}$  for which for all  $i \in G^*$   $x \sim'_i y$  then  $x \sim' y$ .

We are left to show that if  $x \sim' y$  then for all  $i \in G^*$ ,  $x \sim'_i y$ . Assume not. Then there exists  $i \in G^*$  such that  $x \approx'_i y$ . By P we have that  $\exists j \in N - \{i\}$  such that  $x \sim'_j y$ . Let  $G' = \{j \mid x \sim'_i y\}$  so  $i \notin G'$  and  $G' \in \Gamma$ . But if  $G' \in \Gamma$  then  $G^* \subseteq G'$  contradicting  $i \in D$ .

## Question 3

Note that in any equilibrium since the payoff function is linear  $t_i$  satisfies for all i:

$$(*) \quad t_{i} = \begin{cases} 0 & if \ v_{i} < V \\ [0,1] & if \ v_{i} = V \\ 1 & v_{i} > V \end{cases}$$

Since  $v_i$  is monotonically increasing, it follows that if  $t_i > 0$  then then  $t_j = 1$  for all j > i and  $t_j = 0$  for all j < i.

(a)

- For the case where N = 10 and  $v_i = i$  the following is an equilibrium:  $t_i = 0$  for i < 5,  $t_i = 1$  for  $i \ge 5$  and  $V = \sum_{i=5}^{10} \frac{v_i}{10} = 4.5$ . This is an equilibrium. According to (\*)  $t_i$  optimal for all i and  $V = \sum_{i=1}^{N} \frac{t_i v_i}{N}$ .
- For the case where N = 3 and  $v_1 = 1$ ,  $v_2 = 2$  and  $v_3 = 5$  the following is an equilibrium:  $t_1 = 0$ ,  $t_2 = 0.5$ ,  $t_3 = 1$   $V = \frac{0.5v_2 + v_3}{3} = 2$ . Again, according to (\*)  $t_i$  optimal for all i and  $V = \sum_{i=1}^{N} \frac{t_i v_i}{N}$ .

For both cases the equilibrium is unique, since for every V' > V,  $\sum_{i=1}^{N} \frac{t'_i v_i}{N} \leq \sum_{i=1}^{N} \frac{t_i v_i}{N}$  implying that  $V' \neq \sum_{i=1}^{N} \frac{t_i v_i}{N}$  (likewise for V' < V).

(b) Proof of existence of an equilibrium in the general case:

Let N = 1, then  $t_1 = 1$  and  $V = v_1$  is an equilibrium.

Let N > 1. Define  $g(i) = \sum_{j=i}^{N} \frac{v_j}{N}$  and  $v(i) = v_i$  for i = 1, ..., N. (g(i) is the average product per person when only  $j \ge i$  work). Notice that g(i) is strictly decreasing and v(i) is strictly increasing in i, furthermore g(1) > v(1) and g(N) > v(N). So there must exist an  $i^*$  such that  $g(i^* + 1) \le v(i^* + 1)$  and  $g(i^*) > v(i^*)$ . There are two possible cases:  $g(i^*) > g(i^* + 1) \ge v(i^*)$  and  $g(i^*) > v(i^*) > g(i^* + 1)$ .

- For  $g(i^*) > g(i^*+1) \ge v(i^*)$ : let  $V = \sum_{j=i^*+1}^{N} \frac{v_j}{N} = g(i^*+1), t_j = 1$  for all  $j \ge i^* + 1$  and  $t_j = 0$  for all  $j < i^* + 1$ . This is an equilibrium since for all  $j \ge i^* + 1$   $V = g(i^*+1) \le v(i^*+1) \le v(j)$  thus  $t_j = 1$  is optimal and for all  $j < i^* + 1$   $g(i^*+1) \ge v(i^*) \ge v(j)$  thus  $t_j = 0$  is optimal. Finally  $V = \sum_{j=i^*+1}^{N} \frac{v_j}{N} = \sum_{i=1}^{N} \frac{t_i v_i}{N}.$
- For  $g(i^*) > v(i^*) > g(i^*+1)$ : In this case  $i^*$  separates the market into those who do not work  $(j < i^*)$ and the rest who work by choosing  $t^*$  such that  $\sum_{j=i^*+1}^N \frac{v_j}{N} + \frac{t^*v_{i^*}}{N} = v_{i^*}$ . Let  $V = \sum_{j=i^*+1}^N \frac{v_j}{N} + \frac{t^*v_{i^*}}{N} = v_{i^*}$ ,  $t_j = 1$  for all  $j \ge i^* + 1$ ,  $t_{i^*} = t^*$  and  $t_j = 0$  for all  $j < i^*$ . This is an equilibrium since for all  $j \ge i^* + 1$ ,  $V = v_{i^*} < v_j$  thus  $t_j = 1$  is optimal, for  $j = i^* V = v_{i^*}$  thus  $t_{i^*} = t^*$  and for all  $j < i^* V = v_{i^*} > v_j$  thus  $t_j = 0$  is optimal. Finally,  $V = \sum_{j=i^*+1}^N \frac{v_j}{N} + \frac{t^*v_{i^*}}{N} = \sum_{i=1}^N \frac{t_i v_i}{N}$ .